

## Homework 5 Due Mon., Oct. 27

CS 4511: Automata ...

Fall, 2003

**1** Let  $L$  be the language  $\{0^m 1^n \mid m, n \in \mathcal{N}, m \geq n\}$  over  $\{0, 1\}$ . Use the corollary to Theorem 3.2 to show that  $L$  is not accepted by any FA.

**2** Show that there are at least two distinct strings that are indistinguishable wrt the language  $L$  from the previous problem.

**3** Given an FA  $M = (Q, \Sigma, q_0, A, \delta)$ , let

$$Q_r = \{q \in Q \mid \text{there is an } x \in \Sigma^* \text{ s.t. } \delta^*(q_0, x) = q\}.$$

Show that for all  $q \in Q_r$  and  $a \in \Sigma$ ,  $\delta(q, a) \in Q_r$ .

**4** Show that the unreachable states can be eliminated from any FA without affecting the language that it accepts. More precisely, given an FA  $M = (Q, \Sigma, q_0, A, \delta)$ , let  $M_r = (Q_r, \Sigma, q_0, A \cap Q_r, \delta_r)$  where  $Q_r$  is defined as in the previous problem, and for all  $q \in Q_r$  and  $a \in \Sigma$ ,  $\delta_r(q, a) = \delta(q, a)$ . Show that  $L(M) = L(M_r)$ .

(Hint: Begin by showing that for all  $x \in \Sigma^*$ ,  $\delta^*(q_0, x) = \delta_r^*(q_0, x)$ . If your proof uses some kind of induction, specify clearly what the IH is, what you need to show, and where you use the IH.)

**5** Let  $M_1$  and  $M_2$  be the FA's shown below. Draw the FA for  $L(M_1) - L(M_2)$  given by the construction discussed in class. (In light of the result from the previous problem, you may omit any unreachable states.)

**6** For any NFA  $M = (Q, \Sigma, q_0, A, \delta)$ , show that for any  $q \in Q$  and  $a \in \Sigma$ ,  $\delta^*(q, a) = \delta(q, a)$ .

**7** Draw an NFA to recognize  $(a + aa)(ba)^*$ .

**8** Give an example of an NFA  $M = (Q, \Sigma, q_0, A, \delta)$  for which the NFA  $M' = (Q, \Sigma, q_0, Q - A, \delta)$  is such that  $L(M) \cup L(M') \neq \Sigma^*$ . Draw  $M'$ , and give an example of a string not in  $L(M) \cup L(M')$ .

**9** Give an example of an NFA  $M = (Q, \Sigma, q_0, A, \delta)$  for which the NFA  $M' = (Q, \Sigma, q_0, Q - A, \delta)$  is such that  $L(M) \cap L(M') \neq \emptyset$ . Draw  $M'$  and give an example of a string in  $L(M) \cap L(M')$ .

**10** True or False? (In each case, if your answer is false, justify it by giving a counterexample.) For any NFA's  $M_1 = (Q_1, \Sigma, q_1, A_1, \delta_1)$ ,  $M_2 = (Q_2, \Sigma, q_2, A_2, \delta_2)$ , and  $M = (Q_1 \times Q_2, \Sigma, (q_1, q_2), A, \delta)$  where for all  $q \in Q_1$ ,  $r \in Q_2$  and  $a \in \Sigma$ ,

$$\delta((q, r), a) = \delta_1(q, a) \times \delta_2(r, a) :$$

- (a) If  $A = \{(p, q) \in Q_1 \times Q_2 \mid p \in A_1 \text{ or } q \in A_2\}$ , then  $L(M) = L(M_1) \cup L(M_2)$ ?
- (b) If  $A = \{(p, q) \in Q_1 \times Q_2 \mid p \in A_1 \text{ and } q \in A_2\}$ , then  $L(M) = L(M_1) \cap L(M_2)$ ?
- (c) If  $A = \{(p, q) \in Q_1 \times Q_2 \mid p \in A_1 \text{ and } q \notin A_2\}$ , then  $L(M) = L(M_1) - L(M_2)$ ?